

DEFAULT LOGIC

Default Logic is based on inference rules of form

$$\frac{\text{premise } \alpha(x) : \beta(x)}{\delta(x)} \quad \begin{matrix} \nearrow \text{consistency check} \\ \text{(justification)} \end{matrix}$$

(or prerequisite)

E.g., (1) $\frac{\text{Bird}(x) : \text{Flies}(x)}{\text{Flies}(x)}$ Assumes "If it didn't fly, I would know it"!

"Birds generally fly"

"If x is a bird, and it is consistent to assume that x flies, conclude x flies"

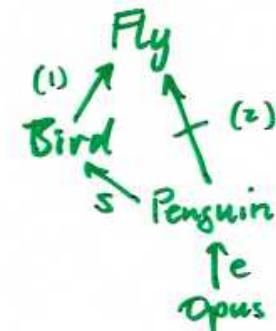
Normal defaults : $\frac{\alpha(x) : \beta(x)}{\beta(x)}$

E.g., add another rule to (1) :

(2) $\frac{\text{Penguin}(x) : \neg \text{Flies}(x)}{\neg \text{Flies}(x)}$

If we also have $\forall x. \text{Penguin}(x) \Rightarrow \neg \text{Bird}(x)$ and $\text{Penguin}(\text{Opus})$, what follows?

cf., default inheritance hierarchies:



see B&L
ch. 10

Two inference paths :

$$\text{Penguin}(\text{Opus}) \vdash_{\text{FOL}} \text{Bird}(\text{Opus}) \vdash_{\text{def. (1)}} \text{Fly}(\text{Opus})$$

$$\text{Penguin}(\text{Opus}) \vdash_{\text{def. (2)}} \neg \text{Fly}(\text{Opus})$$

Note : Once we've applied one of (1), (2), we can no longer apply the other - the consistency test fails!

Compare with CWA, which

- also licences certain conclusions, based on consistency... but disregards new conclusions that may already have been added
- no premises (less "cautious"), only -ve conclusions

We can approximate the effect of CWA with default rules of form

$$\frac{:\neg P(x)}{\neg P(x)}, \quad \frac{:\neg Q(x)}{\neg Q(x)} \quad \text{Note: no premise (or premise = T)}$$

E.g., facts $\{P(A) \vee Q(A)\}$, defaults as above

With CWA, we got an inconsistent augmentation (both $\neg P(A)$ & $\neg P(B)$ are added).

With default logic, once we've added $\neg P(A)$, we can't apply the 2nd default rule, and similarly if we use the 2nd rule first.

"Extensions" of a default theory

Given a db of facts Δ & default rules D , an extension $E[\Delta; D]$ of the default theory is a smallest set s.t.

$$1. \Delta \subseteq E[\Delta; D]$$

2. All logical consequences of $E[\Delta; D]$ are in $E[\Delta; D]$

3. If D contains a rule of form $\alpha(x) : \beta(x)$ & $\alpha(t) \in E[\Delta; D] \wedge \neg \beta(t) \notin E[\Delta; D]$ then $\beta(t) \in E[\Delta; D]$

Inconsistent extensions - can occur for non-normal default theories.

$$\text{E.g., } \{ \}, \frac{\{ :\neg P(A) \}}{\text{facts}}, \frac{:\neg P(A)}{P(A)} \quad \frac{:\neg P(A)}{P(A)} \}$$

Multiple extensions

Normal default theories always have extensions, & these are always consistent - (if Δ is!) (see theorem p. 154). G&N (not mentioned in B&L) However, they can still have multiple extensions.

E.g., earlier example of Opus the penguin: one extension contains $Fly(Opus)$, the other $\neg Fly(Opus)$.

We can modify the rules to get a unique extension:

$$(1') \quad \frac{\text{Bird}(x) : \neg Penguin(x) \wedge Flies(x)}{Flies(x)} \quad \begin{matrix} \text{"semi-} \\ \text{"normal"} \end{matrix}$$

$$\text{or use } (1') \quad \frac{\text{Bird}(x) \wedge \neg Penguin(x) : Flies(x)}{Flies(x)} \quad \begin{matrix} \text{normal,} \\ \text{but requires} \\ \text{more knowledge} \end{matrix}$$

$$(2) \text{ as before } \frac{\text{Penguin}(x) : \neg Flies(x)}{\neg Flies(x)}$$

Given $\text{Penguin}(\text{Opus})$, we can't apply (1') or (1'').

Applying (2) then leads to $\neg \text{Flies}(\text{Opus})$.

$E[\Delta; \delta]$ is then the deductive closure of
 $\{\text{Penguin}(\text{Opus}), \forall x. \text{Penguin}(x) \Rightarrow \text{Bird}(x),$
 $\neg \text{Flies}(\text{Opus})\}$,

i.e., $\{\text{Penguin}(\text{Opus}), \forall x. \text{Penguin}(x) \Rightarrow \text{Bird}(x),$
 $\text{Bird}(\text{Opus}), \neg \text{Flies}(\text{Opus}), \forall x. \text{Bird}(x) \Rightarrow$
 $\text{Bird}(x), \text{Opus} = \text{Opus}, \dots\}$

Semimonotonicity :

Adding new normal defaults to a normal default theory requires no withdrawal of implicit beliefs. (p 154 G&N)

However, the multiple extension problem leads to some peculiar results...

E.5: the Invitation Problem (Kyburg)

"People generally live where their spouses live":

$$\frac{\text{Live-at}(\text{spouse}(x), y) : \text{Live-at}(x, y)}{\text{Live-at}(x, y)}$$

"People generally live where they work":

$$\frac{\text{Work-at}(x, y) : \text{Live-at}(x, y)}{\text{Live-at}(x, y)}$$

Facts: $\text{spouse}(\text{John}) = \text{Mary}$, $\text{Live-at}(\text{Mary}, \text{Buffalo})$,
 $\text{Work-at}(\text{John}, \text{Syracuse})$,
 $\text{Live-at}(x, y) \wedge \text{Live-at}(x, z) \Rightarrow y = z$

Also assume unique names.

We can derive either $\text{Live-at}(\text{John}, \text{Syracuse})$ or
 $\text{Live-at}(\text{John}, \text{Buffalo})$, but (with UNA) not both.
 This is unhelpful if John has invited us and
 we're trying to decide whether to head E
 or W from Rochester!

comment on credulous & skeptical reasoning

E.g. The Lottery Paradox (Kyburg)

1,000,000 tickets have been sold. Generally, a given ticket won't win:

$$\begin{array}{c} \text{Ticket}(x) : \neg \text{Win}(x) \\ \neg \text{Win}(x) \end{array}$$

But some ticket will win:

$$\begin{array}{l} (\exists x) \text{Ticket}(x) \wedge \text{Win}(x) \wedge [x = T_1 \vee \dots \vee x = T_{1,000,000}] ; \\ \text{Ticket}(T_1), \dots, \text{Ticket}(T_{1,000,000}) . \end{array}$$

We can "argue" that $T_{1,000,000}$ will win:

Apply default rule to $T_1, T_2, \dots, T_{999,999}$, obtaining
 $\neg \text{Win}(T_1), \neg \text{Win}(T_2), \dots, \neg \text{Win}(T_{999,999})$. Hence
 (by the existential assertion), $\text{Win}(T_{1,000,000})$!

Similarly we can argue any of the T_i wins,
 i.e., there are 1,000,000 extensions.

(One approach to such problems: allow some arguments to defeat others (e.g., Loui, Pollock))

E.g., Dropouts & Employment - unwanted rule chaining

$$(1) \quad \begin{array}{c} \text{Dropout}(x) : \text{Adult}(x) \\ \text{Adult}(x) \end{array} \quad \begin{array}{l} \text{"Most dropouts are} \\ \text{"adult"} \end{array}$$

$$(2) \quad \begin{array}{c} \text{Adult}(x) : \text{Employed}(x) \\ \text{Employed}(x) \end{array} \quad \begin{array}{l} \text{"Most adults} \\ \text{"are employed"} \end{array}$$

Given that $\text{Dropout}(\text{Joe})$, we conclude
 $\text{Adult}(\text{Joe})$ & hence $\text{Employed}(\text{Joe})$.

Again, we could use more complex versions of (2):

$$(2') \quad \begin{array}{c} \text{Adult}(x) : \neg \text{Dropout}(x) \wedge \text{Employed}(x) \\ \text{Employed}(x) \end{array} \quad \begin{array}{l} \text{semi-} \\ \text{normal} \end{array}$$

or

$$(2'') \quad \begin{array}{c} \text{Adult}(x) \wedge \neg \text{Dropout}(x) : \text{Employed}(x) \\ \text{Employed}(x) \end{array} \quad \begin{array}{l} \text{requires} \\ \text{more} \\ \text{knowledge} \end{array}$$

plus

$$(3) \quad \begin{array}{c} \text{Adult}(x) : \neg \text{Dropout}(x) \\ \neg \text{Dropout}(x) \end{array} \quad \begin{array}{l} \text{or more} \\ \text{rules} \end{array}$$

$$\text{Adult}(\text{Mary}) \xrightarrow{(3)} \neg \text{Dropout}(\text{Mary}) \xrightarrow{(2'')} \text{Employed}(\text{Mary})$$

Other criticisms of default logic :

- as for CWA, drawing default conclusions requires a consistency check, which in general is undecidable.
- default rules enable us to reason with default knowledge, but not about it.

E.g., can't argue as follows:

Quakers are generally pacifists,
pacifists are ^{always} generally Democrats

∴ Quakers are ^{generally} often Democrats

Another approach, "conditional logics", have the opposite problem (defaults are interpretable as facts, but not properly usable as rules)

Probabilistic inference methods can be used to provide a semantic basis for some kinds of default reasoning