CSC2/455 Software Analysis and Improvement Available Expressions, Reaching Definitions, Very Busy Expressions

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URCS

Available Expressions Analysis

The Data Flow Analysis Framework

Reaching Definitions Analysis

Very Busy Expressions (Anticipable expressions)

Outline

Review

Available Expressions Analysis

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Reaching Definitions Analysis

Very Busy Expressions (Anticipable expressions)

- Live variable analysis
 - "Is there a read of this variable along any path?"
- Data flow equations
 - DEF and USE (or UE)
- Iterative data flow analysis

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Very Busy Expressions (Anticipable expressions)

- Identify repeated calculations of the same expression (analysis)
- Remove these repeated calculations (optimization)

a = x * y; if(j == 1) { t = x * y; } else { u = x * y; }

return x * y;



a = x * y; if(j == 1) { t = x * y; } else { x = 3; }

return x * y;



```
Is x * y available at the return
statement?
    a = x * y;
    if(j == 1) {
     x = 5;
     t = x * y;
    } else {
      x = 3;
      u = x * y;
    }
    return x * y;
```



An expression e is available at a program point p iff:

- it has been calculated previously
 - these can be distinct calculations
- all paths from ENTRY to p contain the calculation
- no subexpression of e has been modified along these paths

Informally: Can we reuse the value computed previously by an expression instead of recalculating it?

- What are the similarities from LIVE variable analysis?
- What are the differences from LIVE variable analysis?

- GEN(*n*)
 - a block *n* generates these expressions
- KILL(*n*)
 - a block *n kills* these expressions

- How does information flow for AVAIL?
 - Does a basic block incorporate information from its predecessors?
 - Or its successors?
- How do we "merge" information from multiple source blocks?

 $\text{AVAIL}_{\text{IN}}(n) = ?$

- Information flows forwards
- Only expressions common to *all* paths are available
 - Note: Live variables needed any path

$AVAIL_{IN}(n) = \bigcap_{m \in pred(n)} GEN(m) \cup (AVAIL_{IN}(m) - KILL(m))$

- AVAIL_IN(*ENTRY*) = \emptyset
- AVAIL_IN(n) = {all expressions}

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Building Blocks

- How to generate basic block facts
 - GEN (facts generated by this basic block)
 - KILL (facts killed by this basic block)
- How to merge facts across basic blocks
 - Use a control-flow graph to abstract control flow
 - Direction: forwards/backwards
 - Path requirement: "may" (any) or "must" (all)
 - Determines whether you obtain facts from predecessors or successors, and whether you use union or intersection
 - End goal: a data flow equation for each basic block that needs to be solved
- How to determine initial values
 - Usually analysis specific
- How to solve: iteration until fix point is reached
 - Other methods possible (later in course)

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(Note: The #n notation identifies a definition)

```
y#0 = 0;
x#0 = 1995;
if(y == 0) {
    x#1 = 2000;
    y#1 = 50;
} else {
    x#2 = 3000;
}
a#0 = x + y;
```

Here, the value of y in the last line comes from y#0 or y#1, while the value of x comes from x#1 or x#2.

A definition of variable x at a point p reaches the use of x at a point q iff: there is a path from p to q along which x is not redefined.

- What is the domain of the set of dataflow facts?
- What are the basic block facts (i.e. GEN and KILL)?
- How does information flow?
- How is information merged?

- What is the domain of the set of dataflow facts?
 - The definitions
- What are the basic block facts (i.e. GEN and KILL)?
 - GEN is the set of definitions
 - KILL is the set of definitions
- How does information flow?
 - From a block's predecessors to it
- How is information merged?
 - Union

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Very Busy Expressions (Anticipable expressions)

```
x = 10;
y = 20;
if(x > y) {
    u = x + y;
} else {
    t = x + y;
}
return x + y;
```

- Can you move x + y out of the if to (say) reduce code size?
- Where would you move it?
- What would AVAIL do here?



An expression *e* is *very busy* at the end of a basic block *b* iff:

- *e* is calculated on every path out of *b*,
- *e* could be calculated in *b* without changing the results of calculating *e* in succeeding blocks

Note: *e* may not be currently calculated in *b*. Informally, we're trying to compute an expression earlier than it appears in the code *without* performing a redundant computation.

- What is the domain of the set of dataflow facts?
- What are the basic block facts (i.e. GEN and KILL)?
- How does information flow?
- How is information merged?

- What is the domain of the set of dataflow facts?
 - Expressions
- What are the basic block facts (i.e. GEN and KILL)?
 - GEN is the expressions evaluated in this block
 - KILL is the expressions redefined by this block
- How does information flow?
 - From the block's successors to it
- How is information merged?
 - Intersection

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Very Busy Expressions (Anticipable expressions)

- Chapter 9 of Cooper and Turczon
 - Section 9.2.4
- Also recommended (more detailed coverage here):
 - Aho, Lam, Sethi and Ullman, Chapter 9, Section 9.2.6 (Available Expressions)
 - Aho, Lam, Sethi and Ullman, Chapter 9, Section 9.2.4 (Reaching Definitions)