

CSC2/455 Software Analysis and Improvement Program Analysis

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Outline

And now for something (not so completely) different

Program Analysis in Industry

Basic Notions

Discussion

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Discussion

So far

- ▶ Data flow analysis
- ▶ Loop analysis
- ▶ What next?

Compilers are not the only program analyzers

- ▶ Compilers are probably the most used program analyzers
- ▶ But are severely time constrained
- ▶ Finding *program* errors is not primary goal
 - ▶ Syntax errors, type errors
 - ▶ Code generation primary goal

Program/Software Analysis

- ▶ Software is increasingly mission-critical
- ▶ Can kill people!
 - ▶ Boeing 737 MAX(?)
 - ▶ Therac-25 (X-ray)
 - ▶ Industrial Robotics
- ▶ (less extreme?) Can lose money
 - ▶ Software crashes
 - ▶ Data loss
- ▶ Can we analyze programs for *functional* correctness?
 - ▶ Topic of the next few lectures

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SLAM (Microsoft, early 2000s)

- ▶ MS isolated most crashes to buggy drivers
- ▶ Static Driver Verifier project
 - ▶ Would verify driver code (in C) for correctness
- ▶ Used *model checking*
 - ▶ Models programs as finite-state machines
 - ▶ I used a similar tool (CBMC) to check your assignments



The image shows the word "SLAM" in large, bold, grey 3D block letters. Below the letters, there is a snippet of C code in blue text: `if=node->...; i ++ v1s... end()*node){`. The code is partially obscured by the letters of "SLAM".

Infer (Facebook, early 2010s)

- ▶ Checks C, C++, Objective C, Java and Android code
- ▶ Used for checking Facebook's mobile apps
- ▶ Open source, <https://fbinfer.com/>
 - ▶ Used by Amazon, Mozilla, Uber and Facebook and its affiliates, JD.com, etc.
- ▶ Comes with its own language AL to describe analyses
- ▶ Uses *separation logic*
 - ▶ high-level: converts programs to logic

SPARTA (Facebook, late 2010s)

- ▶ Language-independent analyzer
 - ▶ a C++ framework
- ▶ Open source,
<https://code.fb.com/open-source/sparta/>
- ▶ Used in FB's RedEx tools
 - ▶ for analyzing Android binary code (.dex)
- ▶ Uses *abstract interpretation*
 - ▶ very similar to data flow analysis frameworks



Other efforts

- ▶ Stanford Checker
 - ▶ commercialized by Coverity, late 2000s
 - ▶ CACM article, “A few billion lines of code later: using static analysis to find bugs in the real world”
- ▶ Google’s static analysis tools
 - ▶ Checker Framework for Java programs
 - ▶ Shipshape (abandoned?) (Google Tricorder)
 - ▶ CACM article, “Lessons from Building Static Analysis Tools At Google”
- ▶ Oracle’s Soufflé
 - ▶ Soufflé: Logic Defined Static Analysis

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Limitations

- ▶ None of these frameworks and tools can escape the fact that analysis is an undecidable problem
- ▶ All compute approximations
- ▶ Must be designed to be *sound*
 - ▶ Approximations are conservative/safe
- ▶ Leads to imprecision (i.e. *incomplete*)
 - ▶ May model behaviour not in original programs
 - ▶ (recall IDEAL vs MOP vs MFP)

States and Transitions

- ▶ A program's state is a mapping of variables to values
- ▶ Programs move from one state to another
 - ▶ begin execution in subset of (initial) states
- ▶ Notions of state *before* a program point (i.e. a statement) and *after* a program point
- ▶ Relation that maps before-states to after-states is called a *transition* relation (t)
 - ▶ $\langle x, y \rangle$ (x is before-state, y is after-state)

Traces

- ▶ An execution trace of a program is a sequence of states
 - ▶ $s_0 s_1 s_2 \dots s_n$
- ▶ An execution trace may be finite or infinite
 - ▶ $s_0 s_1 s_2 \dots$
- ▶ The collection of partial traces that can actually happen (i.e. state transitions obey the transition relation) is called the *collecting semantics*
 - ▶ I.e. for all $s_i s_j$ in trace, $\langle s_i, s_j \rangle \in t$

Example

```
x = 0
while(x < 100)
  x = x + 1;
```

- ▶ states are \mathbb{Z}
- ▶ initial state is $\{0\}$
- ▶ transition relation is $\{\langle x, x' \rangle \mid x < 100 \wedge x' = x + 1\}$
- ▶ is 0 1 2 3 part of the collecting semantics?
- ▶ is 0 2 4 6 part of the collecting semantics?

Patrick Cousot and Radhia Cousot, *Basic Concepts of Abstract Interpretation*

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Topics

- ▶ Abstract Interpretation in a Nutshell
- ▶ Abstract Interpretation-based Formal Methods and Future Challenges
- ▶ Liveness Analysis in SPARTA