CSC266 Introduction to Parallel Computing using GPUs Synchronization and Communication

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URCS

Atomics

Warp Primitives

Atomics

Warp Primitives

• Reads and Writes to the same location in shared memory must be separated by a barrier

```
__syncthreads()
```

```
Is this code okay?
```

```
__shared__ volatile int values[8];
int warpid = threadIdx.x / 32;
if(threadIdx.x % 32 == 0)
 values[warpid] = threadIdx.x;
printf("%d\n", values[warpid]);
```

- Does the GPU support global barriers?
- If not, why not?

- Calculate residency R of kernel (use cudaOccupancyMaxActiveBlocksPerMultiprocessor)
- Launch nSM * R blocks, where nSM is number of SMs
- Use global barrier between blocks!
- Usually known as Persistent Threads

- Adds Global Barriers support to CUDA
- And lots more! (See "Cooperative Groups")

```
grid_group grid = this_grid()
grid.sync()
```

- Instead of assuming which blocks are running, discover that at runtime
- Limit synchronization to those blocks

- Volta GPUs no longer execute warps in lockstep
 - "Warp-synchronous" in CUDA literature
 - They can, but it is no longer required
- Every thread has its own PC
- New __syncwarp() barrier for threads in warp
 - Largely for use by code that assumes warp-synchronous execution

Atomics

Warp Primitives

```
atomicCAS(address, compare, val)
```

Pseudocode, ATOMIC means it executes all the instructions inside "atomically".

```
ATOMIC {
   cur = *address;
   if(cur == compare)
        *address = val;
   return cur;
}
```

Other Atomic Functions

- Arithmetic
 - atomicAdd, atomicSub
- Minimum/Maximum
 - atomicMin, atomicMax
- Increase/Decrease
 - atomicInc, atomicDec
 - Always increase or decrease by 1
 - Read definitions carefully!
- Bitwise
 - atomicAnd, atomicOr, atomicXor

- Atomics to same location must be executed serially
- Atomics to different locations in same cache line can be executed in parallel
- Atomics to different locations execute in parallel
 - "As cheap as writes"

• Can we implement all other atomic functions using just CAS?

```
while(atomicCAS(lock, UNLOCKED, LOCKED) == UNLOCKED);
// do something
lock = UNLOCKED;
```

• What happens when two threads in the same warp try to obtain the lock?

Atomics

Warp Primitives

- Allow threads in warps to communicate without using shared memory
- All deprecated in CUDA 9.0 and replaced by more general functions
- See CUDA 8.0 documentation for now

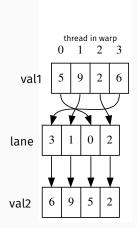
- __any(predicate)
 - Each thread in warp gets 1 if any predicate is 1
- __all(predicate)
 - Each thread in warp gets 1 if only all predicates are 1
- __ballot(predicate)
 - Each thread gets bit pattern of predicates
 - Use __popc() and __ffs for further manipulation

- Allow warps to transfer data to each other
- All functions below are deprecated in CUDA 9.0

T __shfl(T var, int srcLane, int width=warpSize); T __shfl_up(T var, unsigned int delta, int width=warpSize); T __shfl_down(T var, unsigned int delta, int width=warpSize); T __shfl_xor(T var, int laneMask, int width=warpSize);

Warp Permutation using Warp Shuffles

- Assume 4 threads in warp, i.e. 4 lanes
- Can read other lane's values using lane index
- Example: Permutation (pictured)
 val2 = __shfl(val1, lane)
- Example: All lanes read val from lane 0 (broadcast)
 val_t0 = __shfl(val, 0)
- __shfl_up and __shfl_down
 "shift" values across lanes



Adapted from the CUDA programming guide:

```
{
    unsigned int writemask = __ballot(1);
    unsigned int total = __popc(writemask);
    unsigned int prefix = __popc(writemask & __lanemask_lt());
    // Find the lowest-numbered active lane
    int elected_lane = __ffs(writemask) - 1;
    int base_offset = 0;
    if (prefix == 0) {
        base_offset = atomicAdd(p, total);
    }
    base_offset = __shfl(base_offset, elected_lane);
    int thread_offset = prefix + base_offset;
}
```

Atomics

Warp Primitives

```
• Assume all variables below start as 0
```

```
/* Thread 0 */
   val_t0 = 100;
   written_val_t0 = 1;
/* Thread 1 */
```

```
while(write_val_t0 == 0);
printf("%d\n", val_t0);
```

Will the program ever print zero?

In what order are reads and writes in one thread seen by other threads?

- __threadfence_block() X: block
- __threadfence() X: GPU
- __threadfence_system() X: system

All writes made by threads in same X before executing fence are ordered before writes made after executing fence.

```
// writes before fence
fence()
// writes after fence
```